# esproc SPL

SPL

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··· esProc

## Data analysis engine

Low code / High performance / Lightweight / Versatility



## What is esProc?

### What is esProc SPL?

- Data computing and processing engine, run as an analysis **database** or **middleware**.
- Computing and processing of structured and semistructured data
- Offline batch job, online query
- neither SQL system nor NoSQL Technology
- Self created SPL syntax, more concise and efficient



## **SPL: Structured Process Language**

## What pain points does esProc SPL solve?

#### For the data computing scenarios : Offline Batch Job , Online Query/Report



- Slow batch jobs can not fit in certain time window, being strained especially on critical summary dates
- Being forced to wait for minutes for a query/report, the business personnel becomes angry
- More concurrencies, longer the query time span, the database crashes



- ✓ N-layer nested SQL or stored procedures of dozens of KBs, programmer himself is confused after three months
- Dozens of datasources like RDB/NoSQL/File/json/Web/..., cross-source mixed computation is highly needed
- Separate the hot data and cold data into different databases, it is hard to perform real-time queries on the whole data



- Too much relied on the stored procedures, the application can not migrate, framework is hard to adjust
- Too many intermediate tables in the database, exhausting the storage and resources, but dare not to delete them
- Endless report demands in an enterprise, and how can the cost of personnel be relieved?
- .....

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#### What are the counterpart technologies of esProc SPL?

#### Databases that use **SQL** syntax and are applied to **OLAP** scenarios

- Common database: MySQL, PostgreSQL, Oracle, DB2, ...
- Data warehouse on Hadoop: Hive, Spark SQL, ...
- Distributed data warehouse/MPP: ...
- Cloud data warehouse: Snowflake, ...
- All-in-one database machine: ExaData, ...

Other technologies for structured data analysis and statistics

• Python, Scala, Java, Kotlin, ...

#### esProc SPL: Low code, High performance, Lightweight, Openness

## What does esProc SPL bring beyond SQL?



esProc SPL: Reduce the development, hardware, operation and maintenance costs by X times



# **Case Brief**

#### **Case** Batch job of insurance policies of an auto insurance company

T A S K

- Insurance policy table: 35 million rows, details table: 123 million rows
- There are various ways of association, which need to be handled separately



- Informix
- Historical policy matching 6672 seconds
- Codes 1800 lines



- Historical policy matching 1020 seconds
- Codes 500 lines



#### **Case** Batch job of loan agreements of a bank

- T A S K
- SQL: 48 steps, 3300 lines
- Historical data: 110 million rows, daily increase: 1.37 million rows
- Complex multi-table join



- AIX+DB2
- Calculation time: 1.5 hours

esProc SPL

• Calculation time: 10 minutes, codes: 500 lines

8.5 times

**Increase speed** 

Case details: <u>http://c.raqsoft.com/article/1644215913288</u>

#### **Case** Mobile banking: multi concurrent account query

T A S K

Huge number of users, and large concurrent accessesBranch information changes frequently and needs to be

associated in time

- Commercial data warehouses on Hadoop cannot meet the high concurrency requirements
- Using 6 ElasticSearch cluster can cope with concurrency, but can not associate in real time. The data update time is long, and the service must be stopped during this period.

## esProc SPI

- Single machine can cope with the same concurrent volume as ES cluster
- Real time association, zero waiting time for branch information update



**1** server VS **6** servers

#### Case Calculation of the number of unique loan clients of a bank

- Too many labels, and hundreds of labels can be arbitrarily combined to query
- Association, filtering and aggregation calculation of a 20 million rows large table and even larger detailed tables
- Each page involves the calculation of nearly 200 indexes, and 10 concurrency will cause the concurrent calculation of more than 2000 indexes

• Oracle

Α

S

Κ

 Unable to calculate in real time; The query requirements have to be submitted in advance, and the calculation is carried out one day earlier.

## esProc SPL

- 10 concurrency, 2000 indexes in total, less than 3 seconds
- No need to prepare in advance, instantly select any label combination, and get query results in real time

Turn pre-calculation into real-time calculation



## Front-end database in BI System of a bank

Center data warehouse undertakes all data task of whole bank, which is overburdened and can only assign 5 concurrencies to BI system

Only for a small amount of high-frequency data, DB2 is not capable for real-time query, and also unable to achieve data routing, users must select the data source



esProc stores a small amount of high-frequency data, and large low-frequency data is still stored in the data warehouse to avoid repeated construction

esProc takes over the most high frequency computing tasks, and a few low frequency tasks are automatically routed to the center data warehouse

## O An insurance company-Outside database stored procedure

Vertica does not support stored procedures; To prepare data, complex nested SQL statements have to be written, and Java hardcoding is often required.

When mixed computing with MySQL, MySQL data has to be loaded into Vertica first, which is tedious, not real-time, and the database is bloated.

Application layer	Reporting (BIRT)			
Computing layer	esProc-Outside database stored procedures			
Data source	Vertica	MySQL	ACCESS	

#### **User comments**

The best use for us is to pass parameters to the Vertica database.

Each cell becomes a data array that are easy to use, compare and manipulate. It is very logical and you have made it user friendly.



# Why esProc Works Better

# CONTENTS

#### Why SQL is difficult to write: What is the max days has a stock been rising?

SELECT MAX(ContinuousDays) FROM (SELECT COUNT(\*) ContinuousDays FROM (SELECT SUM(UpDownTag) OVER ( ORDER BY TradeDate) NoRisingDays FROM (SELECT TradeDate, CASE WHEN Price>LAG(price) OVER ( ORDER BY TradeDate) THEN 0 ELSE 1 END UpDownTag FROM Stock ) ) GROUP BY NoRisingDays )

SQL doesn't support ordered operation sufficiently and doesn't provide orderly grouping directly; Instead, four layers of nesting has to be used in a roundabout way.

Such statements are not only difficult to write, but also difficult to understand.

In the face of complex business logic, the complexity of SQL will increase sharply, which is difficult to understand and write.

It isn't an unusual requirement, and it appears everywhere in thousands of lines of SQL in reality, which reduces the efficiency of development and maintenance severely.

#### Why can't SQL run fast: Get the top 10 from 100 million rows of data

#### SELECT TOP 10 \* FROM Orders ORDER BY Amount DESC

This query uses ORDER BY. If it is executed strictly according to this logic, it means sorting the full amount of data, and the performance will be poor.

We know that there is a way to perform this operation without full sorting, but SQL cannot describe it. We can only rely on the optimization engine of the database.

In simple cases (such as this statement), many databases can make the optimization, but if the situation is more complex, the database optimization engine will faint

In the following example, get the TopN from each group, SQL cannot describe it directly, and can only write it as a subquery using window function in a roundabout approach.

In the face of this roundabout approach, the database optimization engine cannot do the optimization and can only perform sorting.

SELECT \* FROM (

SELECT \*, ROW\_NUMBER() OVER (PARTITION BY Area ORDER BY Amount DESC) rn

FROM Orders)

WHERE rn<=10

Α

**1** =Stock.sort(TradeDate).group@i(Price<Price[-1]).max(~.len())

The computing logic of this SPL is the same as that of the previous SQL, but SPL provides orderly grouping operation, which is intuitive and concise.

	Α	
1	=file("Orders.ctx").open().cursor()	
2	=A1.groups(;top(10;-Amount))	Top 10 orders
3	=A1.groups(Area;top(10;-Amount))	Top 10 orders of each area

SPL regards TopN as the aggregation operation of returning a set, avoiding full sorting; The syntax is similar in case of whole set or groups, and there is no need to use the roundabout approach.

## Why is SPL more advanced?

【Analogy】Calculate1+2+3+...+100=?

		The smart Gauss came up with a
Ordinary people will do like this	Gauss does like this	convenient and efficient algorithm.
1+2=3	1+100=101	The key here is:
3+3=6 6+4=10	2+99=101 	Multiplication used!
10+5=15 15+6=21 21+7=28	A total of fifty 101 50*101= <mark>5050</mark>	
		Extended reading: SPL: a database language featuring easy writing and fast runnir

SQL is like an arithmetic system with only addition. The code is lengthy and the calculation is inefficient.

SPL is equivalent to the invention of multiplication! Simplify writing and improve performance.

The difficulties of SQL stem from relational algebra, and theoretical problems cannot be solved by engineering methods. Despite years of improvement, it is still difficult to meet complex requirements.

SPL is based on a completely different theoretical system: **discrete dataset**. SPL provides more abundant data types and basic operations, and has more powerful expression capabilities.

## Onter the control of the control

#### **1.** Complex orderly Computing: funnel analysis of user behavior transformation

- Calculate the user churn rate after each event (page browsing, search, shopping cart addition, order placement, payment, etc.)
- Multiple events are effective only when they are completed within a specified time window and occur in a specified order. It is very hard to implement it in SQL, not to mention optimize it.

#### **2.** Multi-step big data batch job

- Complex business requirements are difficult to implement directly in SQL, cursor reading is slow and difficult to calculate in parallel, wasting computing resources
- The implementation using stored procedure requires thousands of lines and tens of steps, with the repeated buffering of intermediate results, the batch job cannot be completed within the specified time window

#### **3.** Multi index calculation on big data, repeated use and multiple associations

- Perform the calculation of hundreds of indexes at one time, and use the detailed data for many times, and association is also involved. SQL needs to traverse data repeatedly.
- Mixed calculation of large table association, conditional filtering, grouping and aggregation, and deduplication; Real-time calculation with high concurrency.

In real business, complex SQL (and stored procedures) are often hundreds/thousands of lines, and a large number of roundabout approaches have to be used to implement the calculation. The code becomes complex as well as the performance becomes low.

## Funnel analysis of an E-commerce company

#### with e1 as ( select uid,1 as step1,min(etime) as t1 from event where etime>= to\_date('2021-01-10') and etime<to\_date('2021-01-25') and eventtype='eventtype1' and ... group by 1), e2 as ( select uid,1 as step2,min(e1.t1) as t1,min(e2.etime) as t2 from event as e2 inner join e1 on e2.uid = e1.uid where e2.etime>= to\_date('2021-01-10') and e2.etime<to\_date('2021-01-25') and e2.etime > t1 and e2.etime < t1 + 7and eventtype='eventtype2' and ... group by 1), e3 as ( select uid,1 as step3,min(e2.t1) as t1,min(e3.etime) as t3 from event as e3 inner join e2 on e3.uid = e2.uid where e3.etime>= to\_date('2021-01-10') and e3.etime<to\_date('2021-01-25') and e3.etime > t2 and e3.etime < t1 + 7and eventtype='eventtype3' and ... group by 1) select sum(step1) as step1, sum(step2) as step2, sum(step3) as step3 from e1 left join e2 on e1.uid = e2.uid left join e3 on e2.uid = e3.uid

	А
1	=["etype1","etype2","etype3"]
2	=file("event.ctx").open()
3	=A2.cursor(id,etime,etype;etime>=date("2021-01-10") && etime <date("2021-01- 25") &amp;&amp; A1.contain(etype) &amp;&amp;)</date("2021-01- 
4	=A3.group(uid).(~.sort(etime))
5	=A4.new(~.select@1(etype==A1(1)):first,~:all).select(first)
6	=A5.(A1.(t=if(#==1,t1=first.etime,if(t,all.select@1(etype==A1.~ && etime>t && etime <t1+7).etime, null))))<="" td=""></t1+7).etime,>
7	=A6.groups(;count(~(1)):STEP1,count(~(2)):STEP2,count(~(3)):STEP3)
SF Cc ef	PL provides order-related calculations and is more thoroughly set-oriented. ode is written directly according to natural thinking, which is simple and ficient.
Th pa	nis code can handle funnels with any number of steps, as long as the nrameters are changed.
SC to	L lacks order-related calculations and is not completely set-oriented. It needs detour into multiple subqueries and repeatedly JOIN. It is difficult to write and

Due to space limitation, only a three-step funnel is listed here, and subqueries need to be added when adding more steps.

understand, and the operation performance is very low.

#### Part of High Performance Computing Mechanism Provided by SPL

Traversal technique	Highly efficient Joins	High performance storage	Cluster computing
Delayed cursor	* Foreign key as pointer	Orderly Compressed Storage	Preemptive Load Balancing
Aggregate Understanding	* Numbering of foreign keys	Free column storage	Multi-zone composite table
		5	
	Order based marge	* Hierarchical Numbering	* Olystar dimension table
* Urdered cursor	Order-based merge	positioning	
Multi-purpose traversal	* Attached table	Index and Caching	Memory spare tire fault tolerance
Prefilter traversal	* Unilateral HASH Join	Double increment segmentation	External storage redundancy fault tolerance

\* Many algorithms and storage schemes here are the original inventions of SPL!

Extended reading: <u>Performance Optimization for Big Structured Data</u>

## **Technical Characteristics**

# CONTENTS

#### esProc Technical Architecture



### Simple and easy-to-use development environment

Bite East Program Tel Window Help         A3 = -create(promo_name,best_sale)         A4 = -create(promo_name,best_sale)         A5 = -create(promo_name,best_sale)         A5 = -create(promo_name,best_sale)         B = -A3 = -create(promo_name,best_sale)         Coloal variable         Cellae previous error message         Copy         Coloal variable         Cellae previous error message         Copy         Coloal variable         Celae previous error message         Copy<		R	Run,	debug, run to cursor, ste	p	Set break	akpoint
A3       = create(promo_name.best_sale)         A3       = create(promo_name.best_sale)         A3       = create(promo_name.best_sale)         A3       = create(promo_name.best_sale)         A4       = stepstore         A5       = stepstore         B4       = stepstore         B4       = stepstore         B5       = stepstore         B4       =		<u>F</u> ile	e <u>E</u> dit	t <u>P</u> rogram <u>T</u> ool <u>W</u> indow <u>H</u> elp			
A3       = create(promo_name,best_sale)         I       = file(*. NdemoVzhVbdWPromotion bt*) import@t()         2       = file(*. NdemoVzhVbdWPromotion bt*) import@t()         3       = create(promo_name,best_sale)         4       3         5       = sale_date<=A4.end_date)         6       = B5 group(clerk_name, - sum(sale_amt):total_amt)         7       = B6 maxp(clotal_amt)         9       = A3         >output(A3)         10       = clean runne         Cellset variable       Clean previous error message         Promo_name best_sale       Feast of St. Fred       Larry         National Pickle Pageant       Steven         0       = message       Copy         8       >output(A3)       Convenient         10       = clean runne       Steven         0       = mercetage       Copy         10       = clean runne       Steven         0       = mercetage       Copy         10       = clean runne       Steven         0       = mercetage       Copy         10       = clean runne       Steven         0       = mercetage       Copy         10       = clea			6	🖬 🕨 🟴 🃭 🛼 👌 S	) II II 🔹 🏟 🚾 🧟 📿 🦻 🗠		
Image: split db11.split       Image: split db11.split db11.spli		A3		= =create(promo_name,best_sale	2)		4
A       B         1       =file(".\\demo\\zh\\bdt\\Promotion.bt").import@t()         2       =file(".\\demo\\zh\\bdt\\SalesRecord L4").import@t()         3       =create(promo_name,best_sale)         4       for A1         5       =A2.select(sale_date>=A4.start_date &&         6       =B5.group(clerk_name,~sum(sale_amt):total_amt)         7       =B6.maxp(total_amt)         8       >A3         9       =A3         >output(A3)    Clean previous error message Copy          promo_name best_sale         Cellest variable             Cellest variable             Real-time system information output, check at any time for abnormality		SPL	db11	splx			
1       =file(*\\demo\\zh\\bdv\\\Promotion bt').import@t()         2       =file(*\\demo\\zh\\bdv\\\SalesRecord bt').import@t()         3       =create(promo_name,best_sale)         4       for A1         5       sale_date<=A4.end_date)         6       =B5.group(clerk_name;~ sum(sale_amt):tota_amt)         7       =B6.maxp(tota_amt)         8       >output(A3)         10       =A3         Simple syntax, in line with natural thinking, simpler than other high-level development languages       Clean previous error message Copy         Promo_name best_sale		\$		А	В		\$ Value
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3       =create(promo_name,best_sale)       1       Feast of SL Fred Larry       What you of the second			2	=file("\\demo\\zh\\txt\\SalesRecor	d.txt").import@t()		Index promo_name best_sale
4       5       for A1       =       A2 select(sale_date>=A4 start_date &&       sale_date<=A4 end_date)			3	=create(promo_name,best_sale)			1 Feast of St. Fred Larry What you see is
5       =A2 select(sale_date>=A4 start_date &&         6       =B5 group(clerk_name;~.sum(sale_amt):total_amt)         7       =B6.maxp(total_amt)         8       >A3         9       =A3         >output(A3)             Simple syntax, in line with natural thinking, simpler than other high-level development languages             Real-time system information output, check at any time for abnormality			4 (	o for A1			2 National Pickle Pageant Steven what you get
6 =B5.group(clerk_name;~.sum(sale_amt):total_amt)   7 =B6.maxp(total_amt)   8 >A3.insert(0,A4.promo_name,B7.clerk_name)   9 =A3   >output(A3)   Simple syntax, in line with natural thinking, simpler than other high-level development languages Real-time system information output, check at any time for abnormality   Real-time for abnormality Task variable			5		=A2.select(sale_date>=A4.start_date && sale_date<=A4.end_date)		cell result ,easy
7       =B6.maxp(total_amt)         8       >A3.inset(0,A4.promo_name,B7.clerk_name)         9       =A3         >output(A3)         10       Convenient cellset variable         Global variable       Watch Output         Cellset variable       Convenient cellset variable         Convenient languages       promo_name best_sale         Real-time system information output, check at any time for abnormality       promo_name best_sale			6		=B5.group(clerk_name;~.sum(sale_amt):tota	al_amt)	to debug;
8       >A3.insert(0,A4.promo_name,B7.clerk_name)         9       =A3         9       =A3         10       >output(A3)         Global variable       Watch Output         Cellset variable       Task variable         Cellset variable       Copy         promo_name best_sale       Feast of St. Fred       Larry         National Pickle Pageant       Steven         Ohristmas Week       Dow			7		=B6.maxp(total_amt)		Convenient to
9 =A3     10 </th <th></th> <th></th> <th>8</th> <th></th> <th>&gt;A3.insert(0,A4.promo_name,B7.clerk_name</th> <th>e)</th> <th>reference</th>			8		>A3.insert(0,A4.promo_name,B7.clerk_name	e)	reference
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Real-time system information output, check at any time for abnormality				developm	entianguages		promo_name best_sale
				Real-time syste any t	m information output, check a	t	Feast of St. Fred     Larry       National Pickle Pageant     Steven       Christmas Week     Dow

## Specially designed syntax system

#### SPL is especially suitable for complex process operations.

	А	В	С	D	E	F
1	=esProc.query("SELECT or	derID as contract, orderDate	e as date, customer, amour	t, emplD as salesman FRO	M sales where year(orderD	ate)=? OR year(orderDate)
2	=esProc.query(select * fron	n employeeInfo")				
3	≻A1.run(salesman=A2.sele	ect@1(ID:A1.salesman))	/field value is record			
4	≻A1.group(salesman)					
5	=create(salesman, thisyea	rAmount, lastyearAmount, g	rowthRate, custNumber, big	CustNumber,bigCustPropo	rtion)	
6 😑	for A4	=A6(1).salesman.name				
7		=A6.select(year(date)==yea	r).sum(amount)			
8		=A6.select(year(date)==yea	r-1).sum(amount)			
9		=B8/B7-1	/growth rate			
10		=A6.group(customer).(~.su	m(amount))			
11		=B10.count()	/number of customer			
12		=B10.count(~>=10000)	/number of big customer			
13		=B12/B11				
14		=A5.insert(0,B6,B7,B8,B9,B	11,B12,B13)			

Natural & clean step-by-step computation, direct reference of cell name without specifically defining a variable

### Rich class libraries

#### Designed specifically for structured data tables

	A	В	С	
1	=esProc.query("SELECT or	rderID as contract,	/retrieve sales records	
2	=A1.group(salesman)			
3	=create(salesman,thisyear	Amount, lastyearAmount, cu	istNumber, bigCustNumber	)
4 G	for A2	=A4(1).salesman		
5		=A4.select(year(date)==yea	ir).sum(amount)	
6		=A4.select(year(date)==yea	ar-1).sum(amount)	
7		=A4.group(customer).(~.su	m(amount))	
8		=B7.count()		

#### **Grouping & Loop**

	А	В	С
1	=esProc.query("select * fro	m employee")	
2	=A1.select(sex=="male")		
3	=A1.select(birthday≻=date(	'1970-01-01"))	
4	=A2^A3	/intersect, find out male em	ployee born after 1970
5	=A2&A3	/union, find out-male empl	oyee or employee born after 1
6	=A2VA3	/subtract, find out male emp	oloyee born before 1970
7	=A4.sum(salary)		
8	=A5.avg(age)		
9	=A5.sort(birthday)		

#### Set operations

	A	В	С
1	=file("traderecord.txt").impo	rt@t()	
2	=A1.sort(customerID, trade	Date)	
3	=A2.select(autoType=="Jet	a"    autoType="Passat").du	p@t()
4	=A3.derive(interval(tradeDa	te[-1], tradeDate):space)	
5	=A4.select(autoType[-1]=="	Jetta" &&atuoType=="Passa	t" &&customerID=customerI
6	=A5.avg(space)		
7			
0			

#### **Sorting & Filtering**

	А	В	С
1	=esProc.query("select * from employee")		
2	=A1.sort(entryDate)		
3	=A2.pmin(birthday)	/select recordNo of employ	ee born at earl
4	=A2(to(A3-1))	/directly access employee r	ecord via recol
5	=esProc.query("select * from stock where	stockCode="000062")	
6	=A5.sort(tradeDate)		
7	=A6.pmax(closePrice)	/recordNo of highest excha	nge closing qu
8	=A6.calc(A7,closePrice/closePrice[-1]-1)		

#### **Ordered sets**

## Excellent integration

esProc, developed in Java, can run independently or be seamlessly integrated into applications.



## Support diverse data sources

Multiple data sources are directly used for mixed computations.



Direct computation without loading data into database, using the strengths of original data sources, real time computation

#### Flexible and efficient file storage

#### High performance

Private data storage format: bin files, composite tables

#### File system

Support to store data by business classification in tree directory

#### BIN FILE

COMPOSITE TABLE Double increment segmentation supports arbitrary number parallelism Self-owned efficient compression technique (reduced space; less CPU usage; secure) Generic Storage, allowing set data

#### Mixed row and column storage

Ordered storage improves compression rate and positioning performance

#### Efficient intelligent index

**Double increment segmentation** supports arbitrary number parallelism **Integration of main and sub table** to reduce storage and association **Numbering keys** to achieve efficient positioning Join

Direct file storage without database is more efficient, more flexible, and cheaper.



# **More Solutions**

## Implementation of data-driven micro-service

• Mainstream frameworks such as microservices require data processing at the application side.

#### Problems

- The database is difficult to be embedded in front-end applications, thus hardcoding is the only choice.
- Java/ORM lacks sufficient structured computing class library, which makes it difficult to develop data processing, and hot swap can not be achieved.



- SPL replaces Java/ORM to implement data computing in (micro-service) applications.
- Rich class library and agile syntax simplify the development.
- ✓ The system is open and can process data of any source in real time.
- SPL is interpreted executed, naturally supporting hot swap.
- Efficient algorithms and parallel mechanism ensure computing performance.

## **Replace stored procedures**

•	Stored procedures	are hard to edi	t and debug,	and lack mig	gratability.
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#### Problems

- Compiling stored procedures requires high privilege, causing poor security. ٠
  - The shared use of a stored procedure by multiple applications will cause tight coupling between applications. ٠



- SPL is intuitively suitable for complex multi-step data
- SPL scripts are naturally migratable.
- The script only requires the read privilege of the database and will not cause database security problems.
- Scripts of different applications are stored in different directories, which will not cause coupling between

## Eliminate intermediate tables from databases

• For query efficiency or simplified development, a large number of intermediate tables are generated in the database.

#### Problems

- The intermediate tables take up large space, causing the database to be excessively redundant and bloated.
- The use of the same intermediate table by different applications will cause tight coupling, and it is difficult to manage the intermediate tables (hard to delete).



- The aim for storing intermediate tables in the database is to employ the database's computational ability for subsequent computations; SPL can implement the subsequent computations after using file storage.
- External intermediate tables (files) are easier to manage, and using different directories for storage will not cause coupling problems between applications.
- External intermediate tables can fully reduce the load on the database, even without the need to deploy the database.

## Distance in the second seco

	•	Reporting tools/BI tools can only solve the problems in the report presentation stage and can do nothing about data preparation.
Problems	•	Data preparation implemented in SQL/stored procedure/Java hardcoding is difficult to develop and maintain, and the cost is high.

• The report development needs are objectively endless, and data preparation is the main factor leading to high development costs.



- Add a computing layer between report presentation and data source to solve the data preparation problems.
- SPL simplifies the data preparation of reports, makes up for the lack of computing ability of reporting tools, and comprehensively improves the efficiency of report development.
- Both report presentation and data preparation can quickly respond to handle endless report development needs at low cost.

## Programmable data routing to implement front-end calculation

Center data warehouse needs front-end calculation to share the heavy workload

- Move only high frequency data to front, take over most of the calculation requests
- The programmable data routing automatically selects the front-end data and data warehouse and mixes the calculation results, applications have transparent access to full data



Data routing enables low-cost high-performance full data analysis

- Dilemmas of implementing front-end calculation through traditional database
  - Application may access all the data, and moving full data to front causes redundant construction and high cost
  - Due to lack of data routing, moving part of the data to front causes that application can not transparently access the full data and the poor experience

#### Mixed computation to implement real-time HTAP

- Mixed calculation of hot and cold data to implement T+0 real-time analysis
  - Organized historical cold data
  - Real time reading of transaction hot data
- The production system needs no modification
  - Make full use of the advantages of the original multiple data sources



Support low-risk, high-performance and strong real-time HTAP with open multi-source hybrid computing capability

- HTAP databases are difficult to meet HTAP requirements
  - Requiring replacement of production system database, with high risk
  - Insufficient SQL computing power, insufficient historical data preparation, low performance
  - Closed computing capacity, complex ETL process required due to external diverse data sources, poor real-time ability

## Perform computation on files to implement Lakehouse

- Open format file data computation
  - txt/csv/xls/json/xml
- High performance private format file storage and computation
- Enter at will and sort out gradually; Lake to House
- Rich data source interfaces, direct real-time computing



Data Source

Data Lake

- RDB can only House, not Lake
- Strong constraints, non-compliant data cannot enter, and complex ETL processes are inefficient
- Closed, external data cannot be calculated, let alone mixed real-time calculation

# **Frequently asked questions**

#### Is esProc based on open source or database technology?

esProc is based on a brand-new computing model, no open source technology can be cited, and all independent innovation from theory to code.



SPL is based on innovation theory that can no longer use SQL to achieve high performance, and SQL can not describe most low complexity algorithms.

## **Q** Where can esProc be deployed

esProc is implemented in pure Java.

esProc can run smoothly under any OS equipped with JVM, including VM, cloud server and even container.

## How applications invoke esProc?

esProc provides a standard JDBC driver for Java applications.

esProc can be integrated in a Java application seamlessly.

esProc can be invoked by a non-Java application via HTTP/RESTFul

## Can esProc run based on the existing database?

#### Yes, Of course!

However, esProc can not guarantee high performance in this situation due to the inefficient I/O of database, and database can hardly provide storage schema which is necessary for low complexity algorithm.

## Where does esProc store data?

esProc stores data in files of private format to guarantee performance.

Any file system installed on OS is available, including NFS.

esProc can easily implement separation between storage and computation to scale out



## What are the weaknesses of esProc?

#### **Comparing with RDB:**

The metadata is immature in esProc, most of computation will begin from accessing data source, it will be a little tedious for simple operations.

#### **Comparing with Hadoop/MPP:**

The cluster function of esProc has not many chances to be well-trained. esProc has reduced many clusters into a single machine without sacrificing performance in history.

#### **Comparing with Python:**

SPL is developing its AI functions, but now is still not even close to Python.



SPL is dedicated to low code and high performance.

Learning SPL is not difficult. It can be mastered in hours and skilled in weeks. What is difficult is to design optimization algorithms!



## Is there a tool to convert SQL to SPL automatically?

Not yet.

The information in SQL statement is insufficient to optimize its performance. Frankly, we are not a veteran like RDB vendor for guessing goal of a SQL, so converting SQL to SPL directly will usually lead to slower speed.



## O How to launch a performance optimization process

#### The first 2-3 scenarios will be implemented by Scudata engineer in collaboration with users.

Most programmers are used to the way of thinking in SQL and are not familiar with high performance algorithms. They need to be trained to understand in one or two scenarios.

Dozens of performance optimization routines will be experienced and learned. Algorithmic design and implementation are not so difficult.



Give a man a fish and you feed him for a day. Teach him how to fish and you feed him for a lifetime!

Extended reading : <u>What should we do when SQL (and stored procedure) runs too slowly?</u>

#### Summary of advantages of esProc



#### **High performance**

The processing speed of big data is 1 order of magnitude higher than that of traditional solutions



#### **Efficient development**

Procedural syntax, in line with natural thinking

Rich class libraries





#### Flexible and open

Multi-source mixed computation Can run independently, or embedded into applications

#### Save resources

Single machine can match cluster, reducing hardware expense Environment-friendly



#### Sharp cost reduction

Development, hardware, O&M costs reduced by X times

# THANKS